

# SPONGENT: A Lightweight Hash Function

**Andrey Bogdanov**, Miroslav Knežević, Gregor Leander,  
Deniz Toz, Kerem Varıcı, Ingrid Verbauwhede  
K.U.Leuven, NXP, DTU

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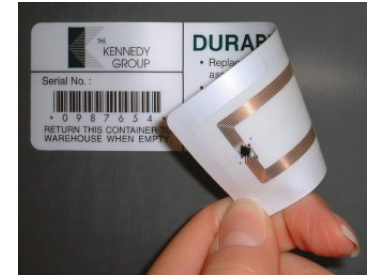
# Ubiquitous Computing



Implants



Sensor networks



Logistics



Transportation



Access control



IDs

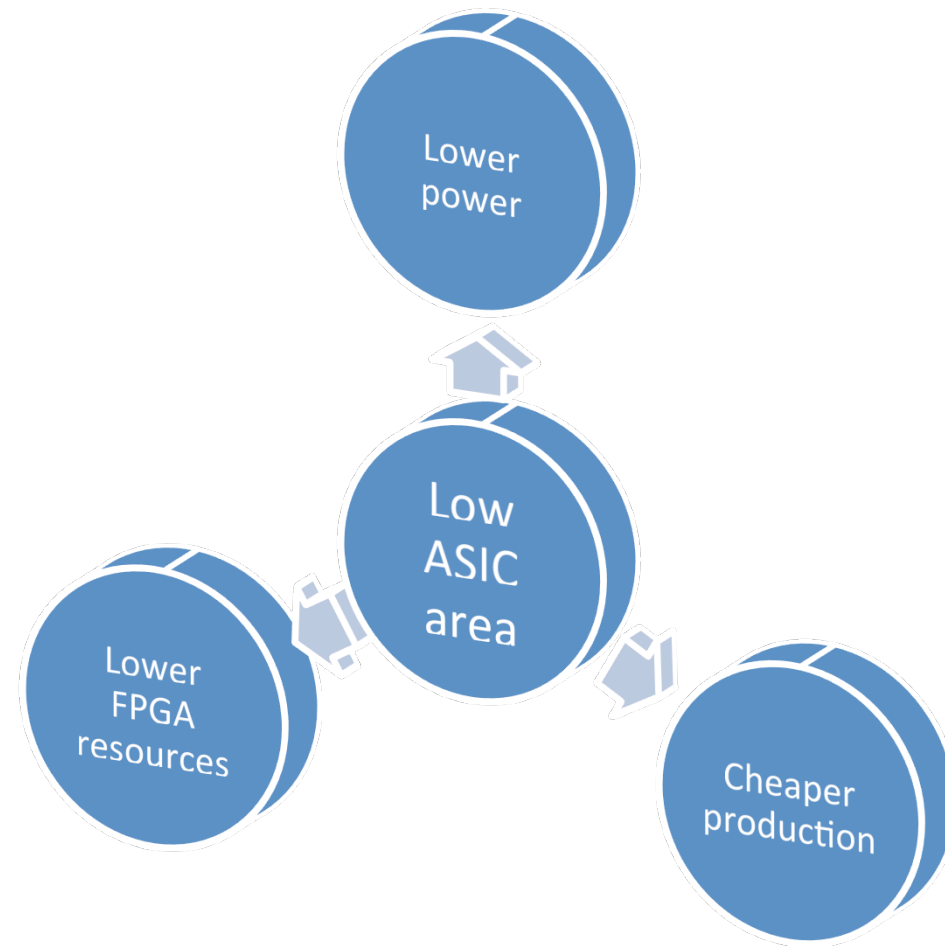
# Lightweight = Low Cost

Lightweight = low development costs?  
Lightweight = small code size?

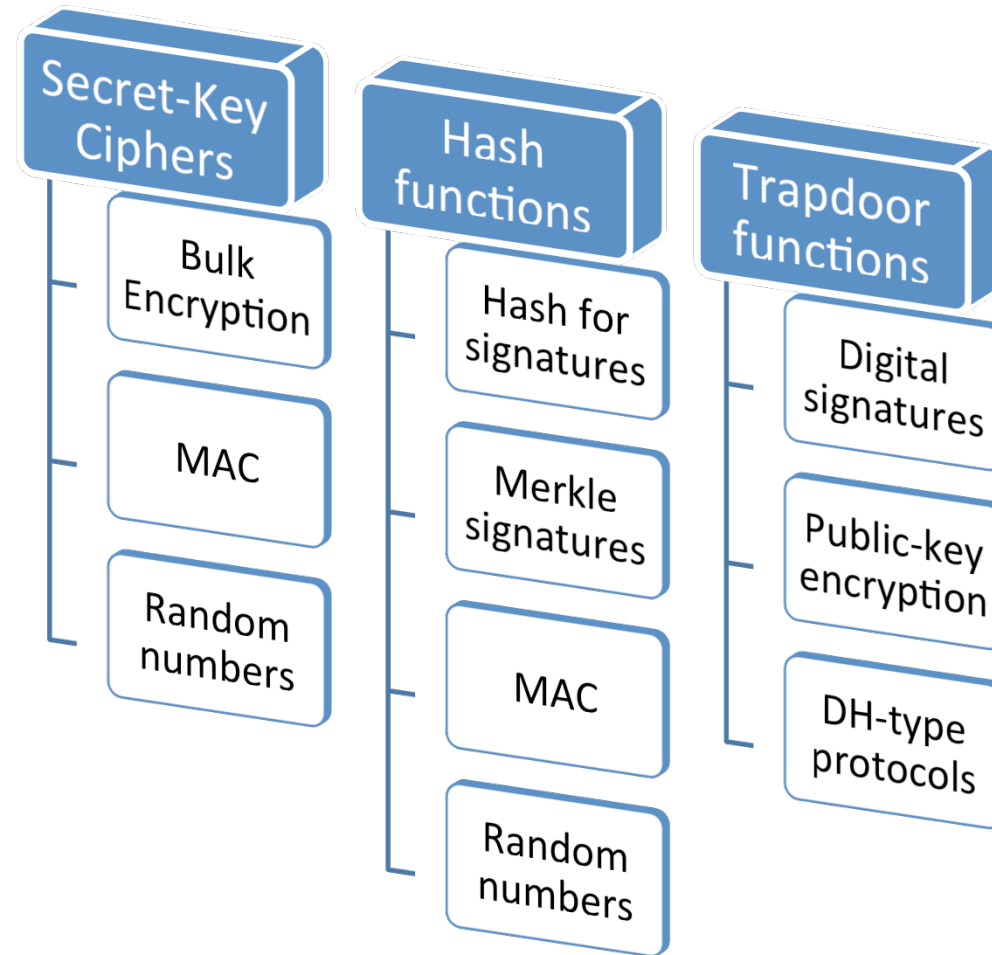
Lightweight = low ASIC area?  
Lightweight = low usage of FPGA resources?

Lightweight = low power consumption?  
Lightweight = low energy consumption?

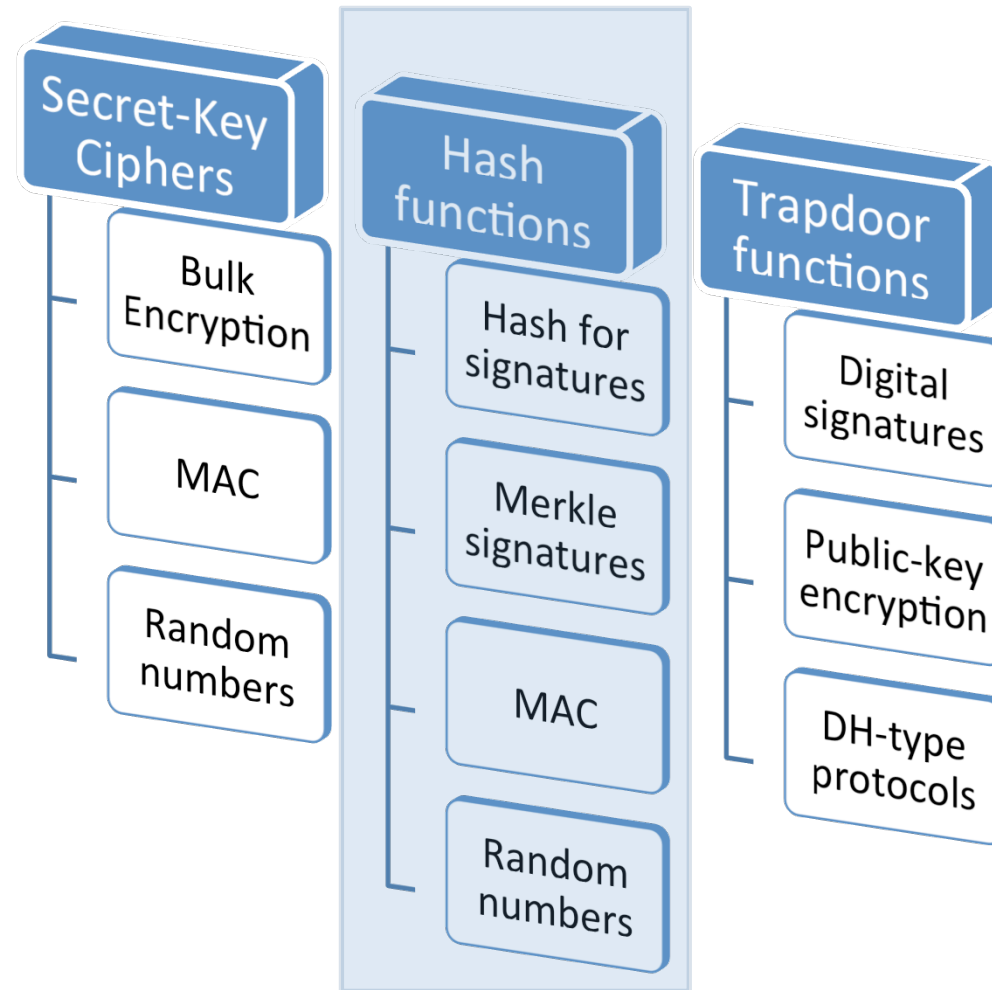
# Lightweight = Low ASIC Area



# Crypto Algorithms to Meet Basic Security Needs

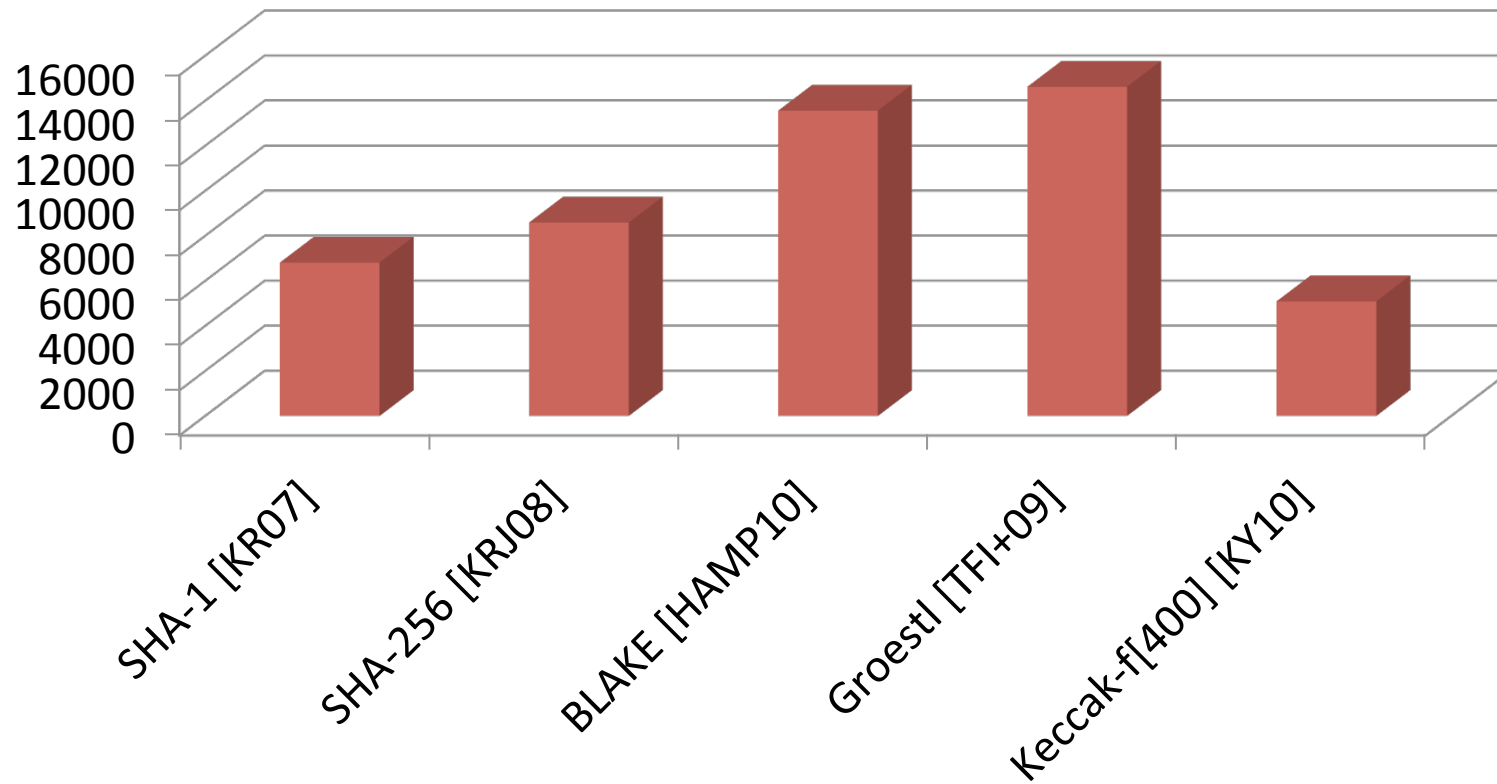


# Crypto Algorithms to Meet Basic Security Needs



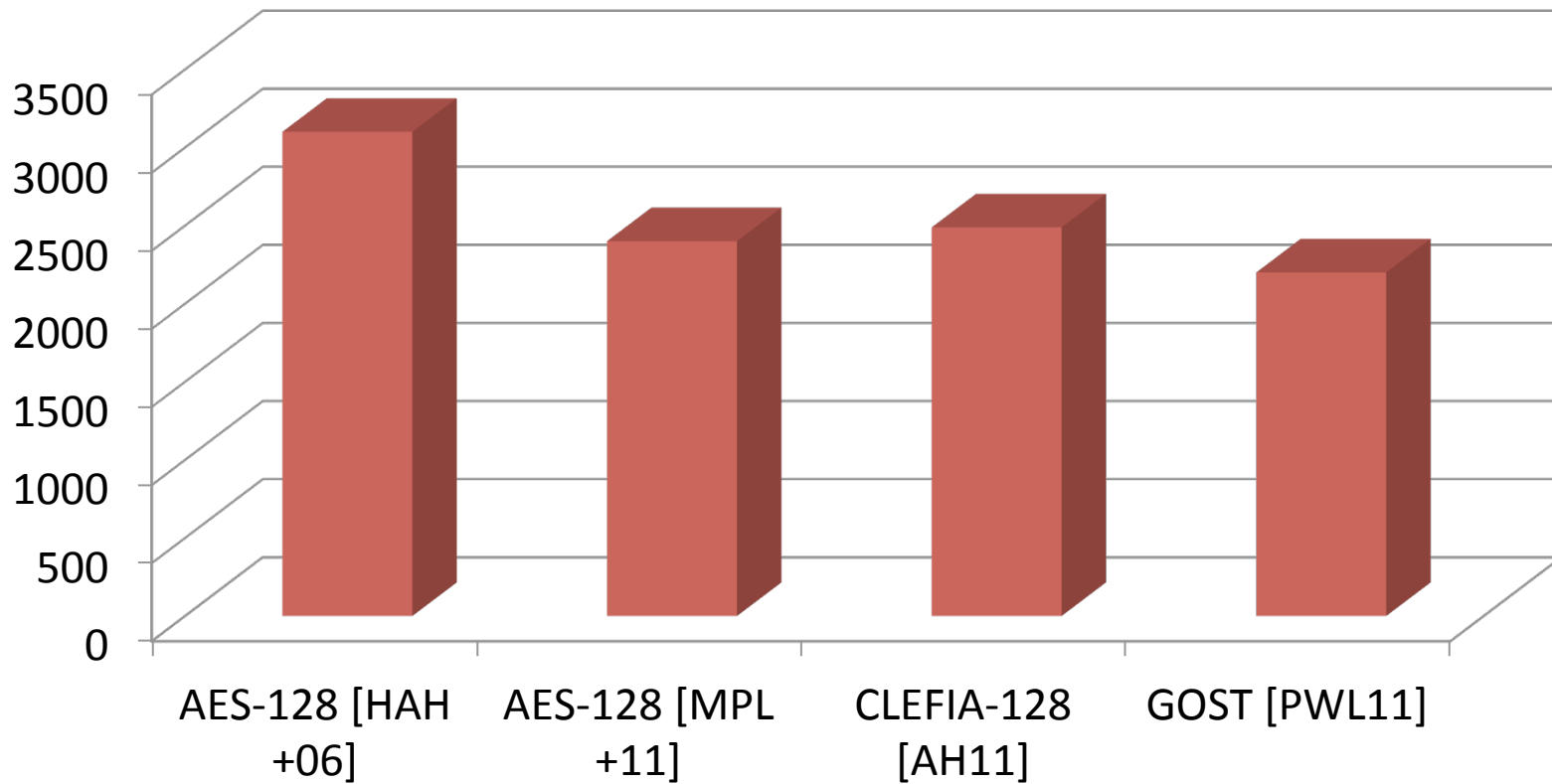
# General-Purpose Hash Functions

Area in Gate Equivalents (NAND gates)



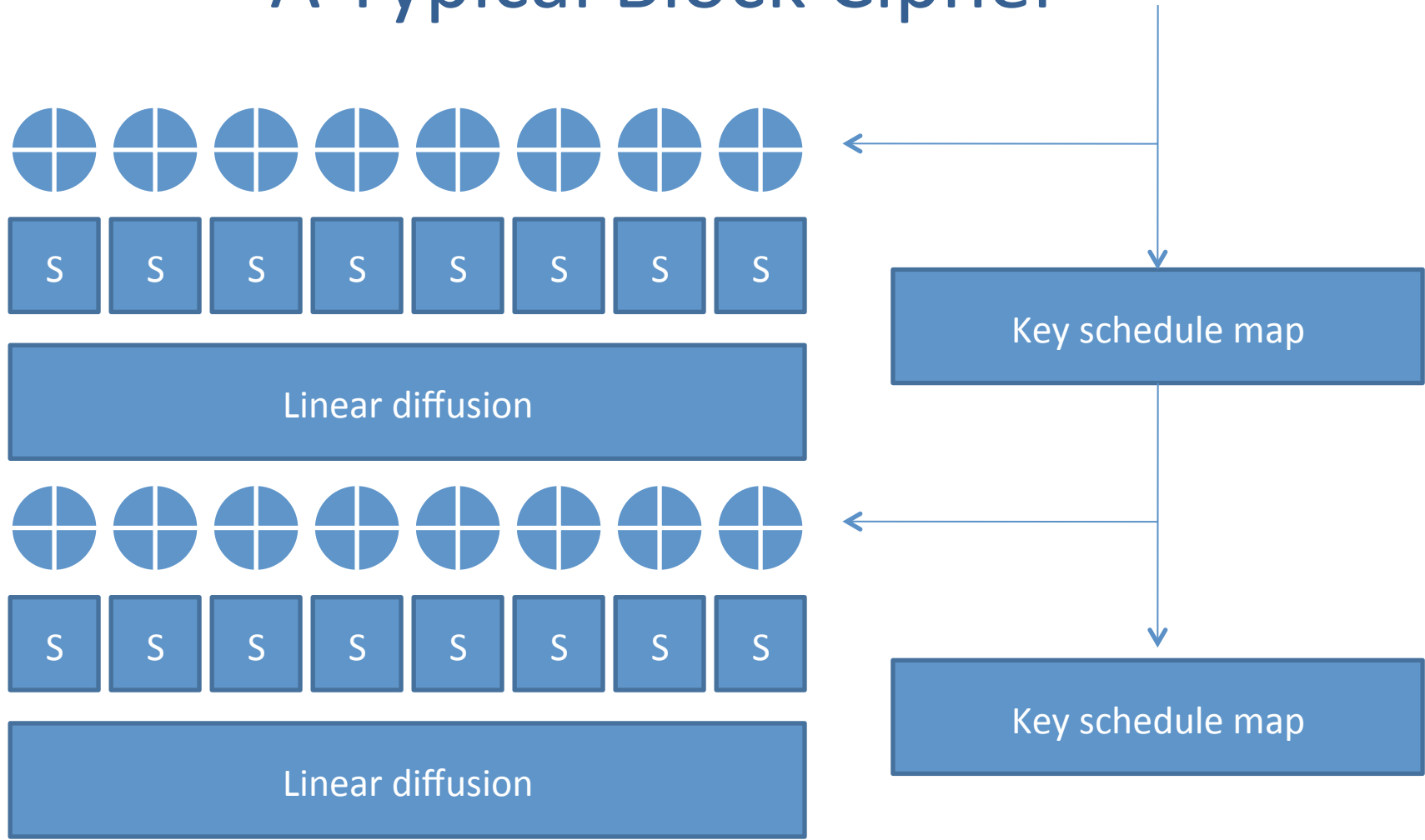
# Some Block Ciphers

Area in Gate Equivalents (NAND gates)



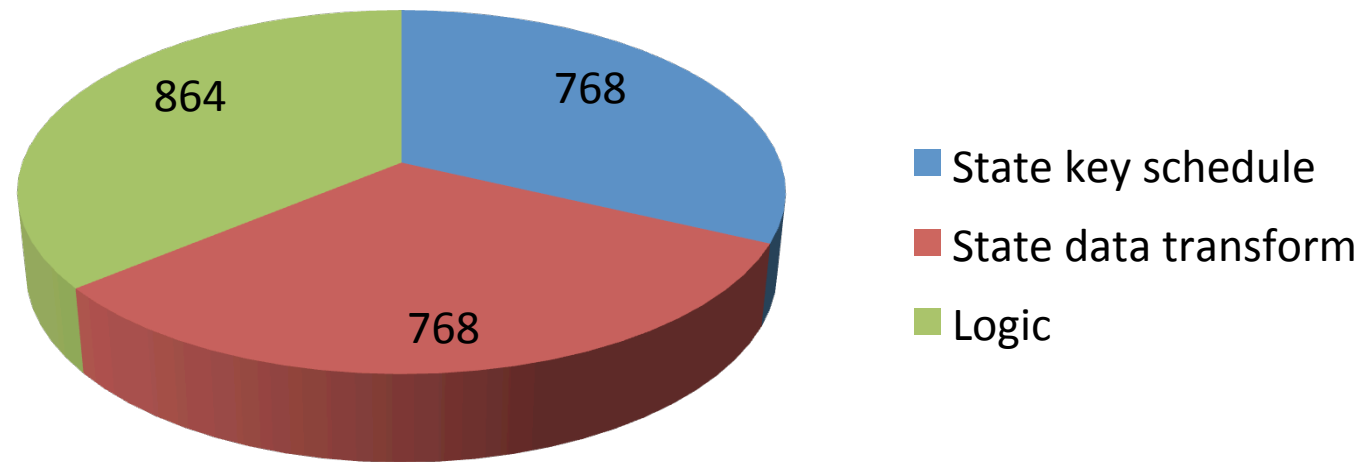


# Substitution-Permutation: A Typical Block Cipher



# Contribution of Building Blocks

Approximate area of AES-128 [MPL+11] in GE



**How to optimize the design of a cipher for low area?**  
For given block and key sizes, minimize logic!

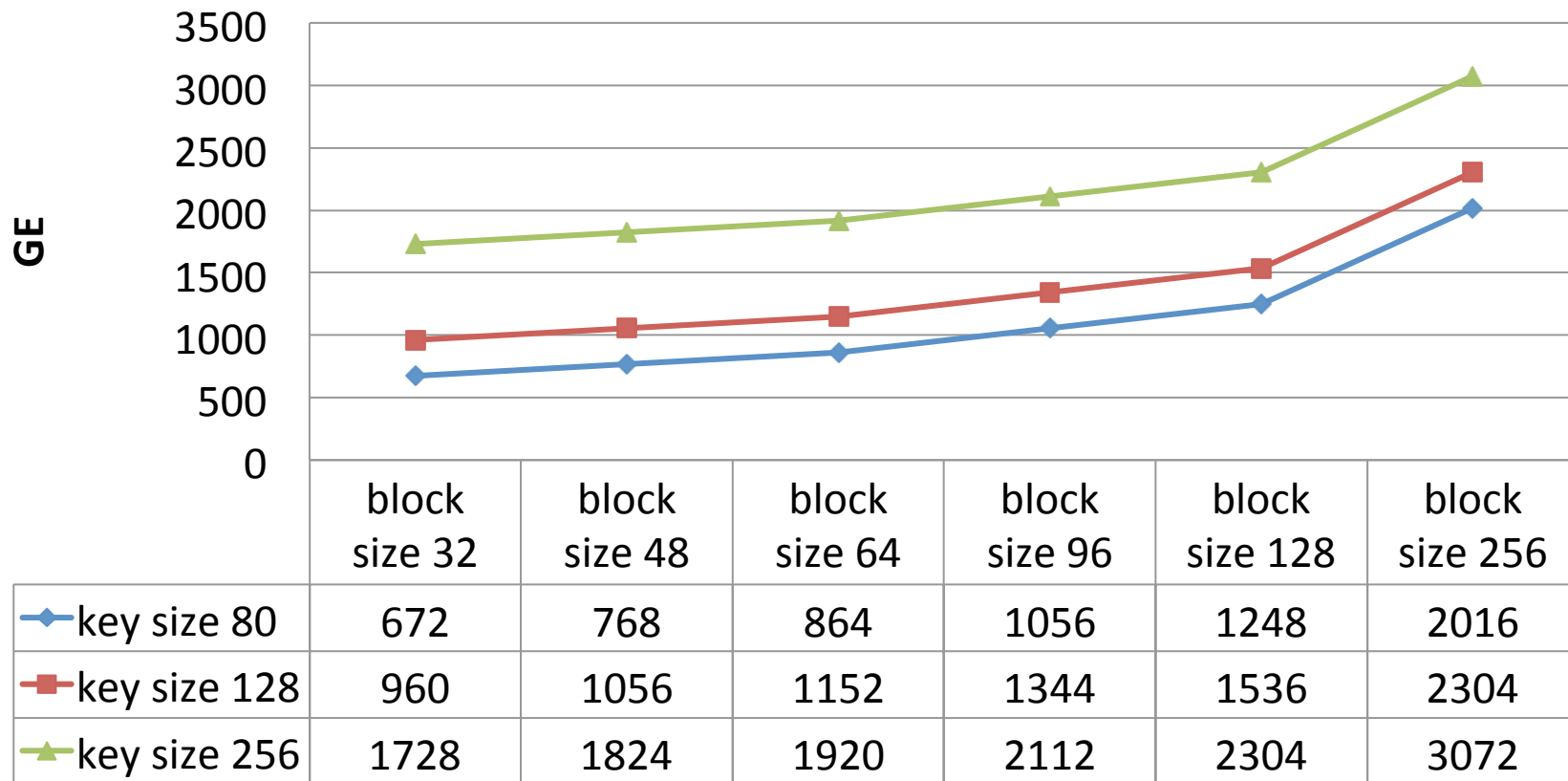
# Area of Elementary Blocks

Logical operation	Cost in ASIC hardware
<b>NAND(x,y)</b>	<b>1.00 GE</b>
AND(x,y)	1.25 GE
OR(x,y)	1.25 GE
XOR(x,y)	2.25 GE
MUX(x,y;c)	2.50 GE
AND(x,y,z)	1.50 GE
MAJ(x,y,z)	2.25 GE
<b>XOR(x,y,z)</b>	<b>4.00 GE</b>

State (flip-flop)	Cost in ASIC hardware
1 bit	5.50 -7.50 GE

# Minimum Area for Security Parameters

Lower bounds on area for block ciphers (for 1 FF = 6GE)



# Design Ideas

**Small State:** Minimize block and key sizes

Each bit removed saves about 5 GE

Tailor sizes to your exact needs



**Small Logic:** Minimize algorithm description

Min nonlinear = 1 GE per NAND

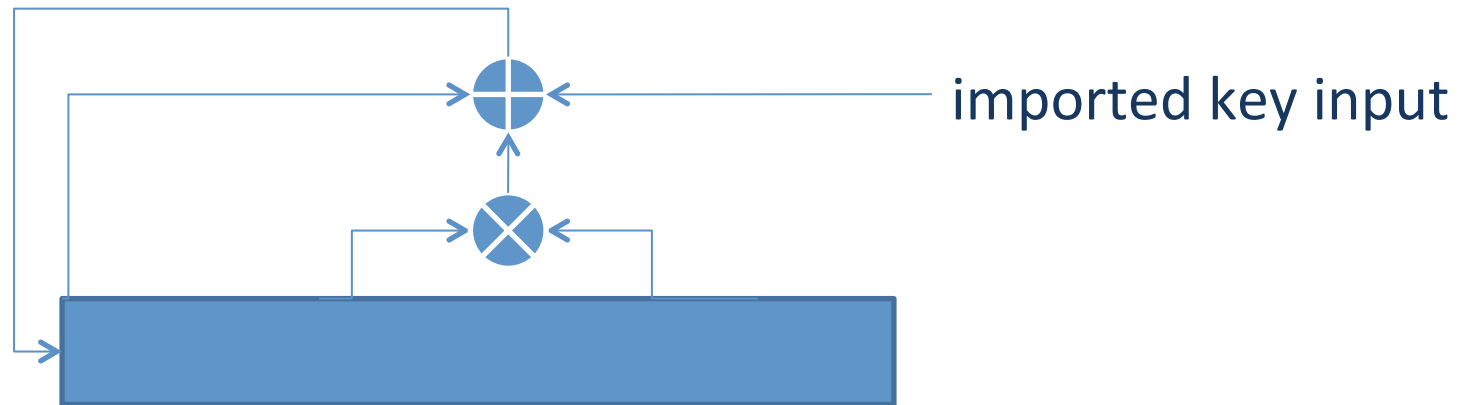
Avoid linear = about 2 GE per XOR



**No key schedule:** Get rid of the on-the-fly key schedule

Can save up to 50% or more area

# The Extreme Lightweight Cipher



n-bit block

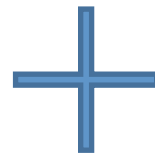
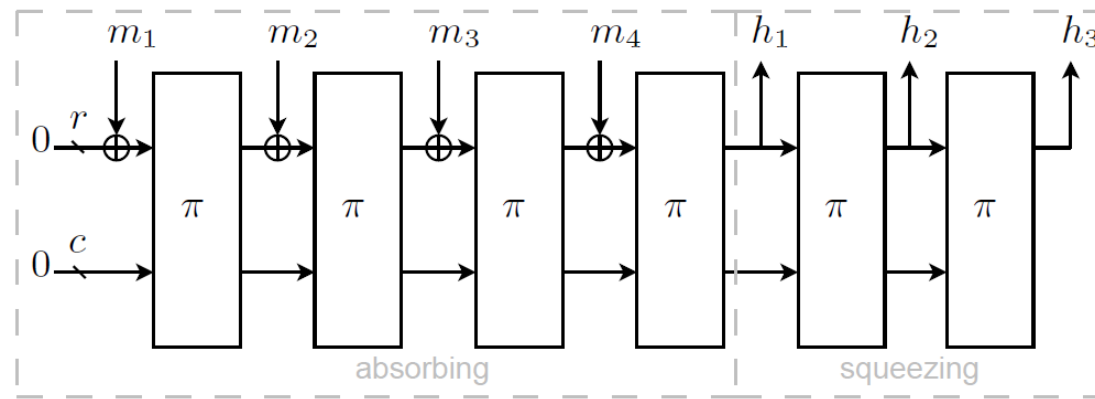
At least several thousand rounds required to attain security!

block size n	area: 5+6n GE
32	192
48	293
64	384
96	581
128	773

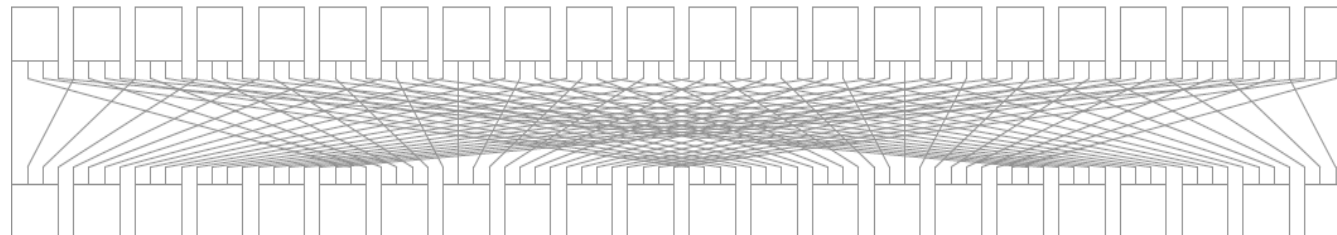
# SPONGENT Hash: The Overview



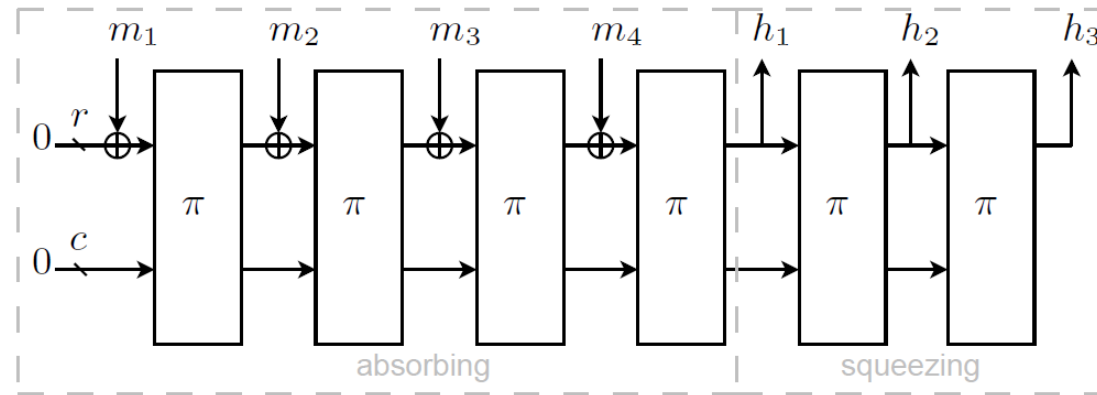
SPONGE construction



Unkeyed **PRESENT**-type permutation  $\pi$ : 4-bit S-box and bit diffusion



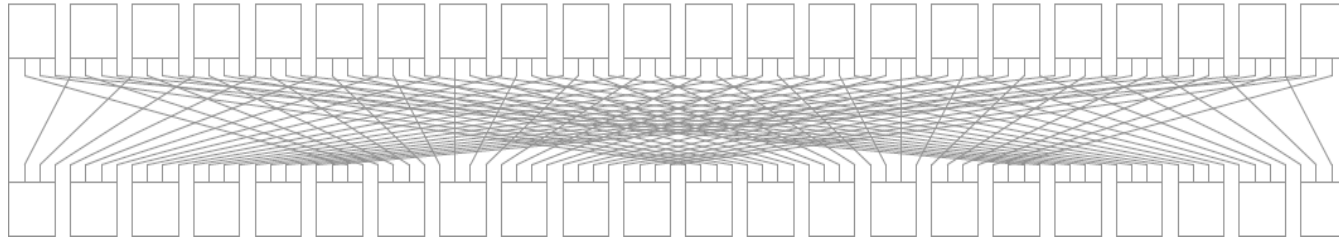
# SPONGENT: The Parameters



	$n$ (bit)	$b$ (bit)	$c$ (bit)	$r$ (bit)	$R$ number of rounds	security(bit)		
						preimage	2nd preimage	collision
SPONGENT-88	88	88	80	8	45	80	40	40
SPONGENT-128	128	136	128	8	70	120	64	64
SPONGENT-160	160	176	160	16	90	144	80	80
SPONGENT-224	224	240	224	16	120	208	112	112
SPONGENT-256	256	272	256	16	140	240	128	128



# SPONGENT: Building Blocks



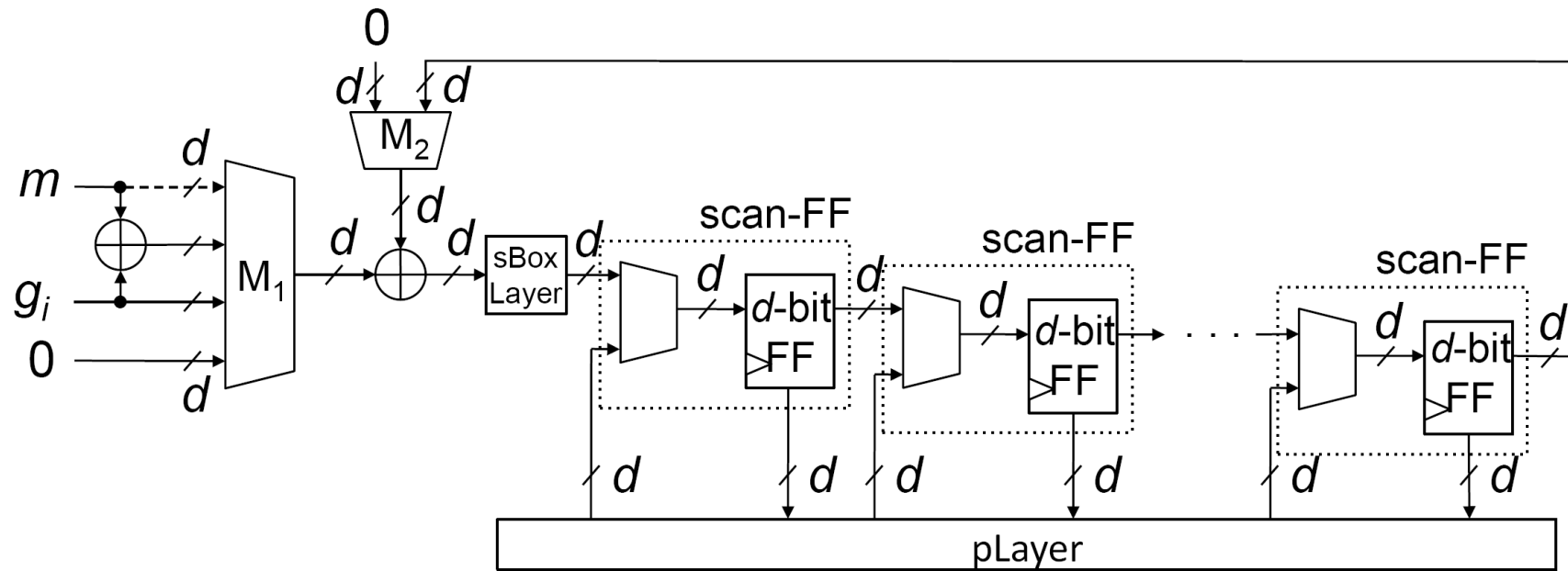
**New 4-bit S-box fulfilling the PRESENT S-box criteria:**

$x$	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
$S[x]$	E	D	B	0	2	1	4	F	7	A	8	5	9	C	3	6

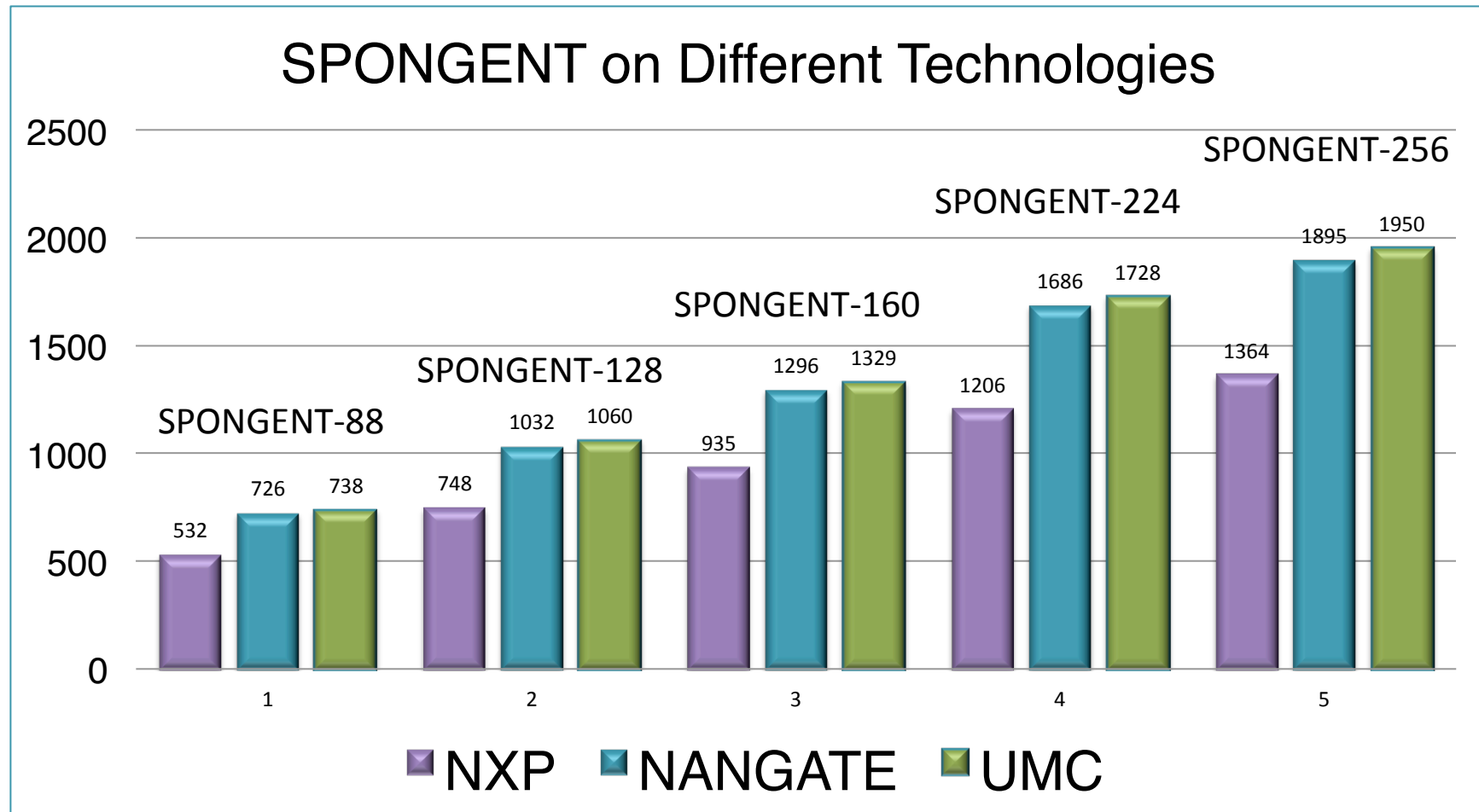
**PRESENT-type bit permutation:**

$$P_b(j) = \begin{cases} j \cdot b/4 \bmod b - 1, & \text{if } j \in \{0, \dots, b - 2\} \\ b - 1, & \text{if } j = b - 1. \end{cases}$$

# SPONGENT: Serial ASIC Implementation



# SPONGENT: The Impact of Technology



# SPONGENT: The Area, UMC 130nm

Hash Function	Cycles	Datapath [bit]	Area [GE]	Throughput [kb/s]
SPONGENT-88	900	4	<b>738</b>	0.81
	45	88	<b>1127</b>	17.78
SPONGENT-128	2380	4	<b>1060</b>	0.34
	70	136	<b>1687</b>	11.43
SPONGENT-160	3960	4	<b>1329</b>	0.40
	90	176	<b>2190</b>	17.78
SPONGENT-224	7200	4	<b>1728</b>	0.22
	120	240	<b>2903</b>	13.33
SPONGENT-256	9520	4	<b>1950</b>	0.17
	140	272	<b>3281</b>	11.43

# SPONGENT: The Area, NANGATE 45 nm

Hash Function	Cycles	Datapat h [bit]	Area [GE]	Throughput [kb/s]	Area compared to UMC
SPONGENT-88	900	4	<b>726</b>	0.81	<b>-2%</b>
	45	88	<b>1164</b>	17.78	<b>+3%</b>
SPONGENT-128	2380	4	<b>1032</b>	0.34	<b>-3%</b>
	70	136	<b>1706</b>	11.43	<b>+1%</b>
SPONGENT-160	3960	4	<b>1296</b>	0.40	<b>-2%</b>
	90	176	<b>2272</b>	17.78	<b>+4%</b>
SPONGENT-224	7200	4	<b>1686</b>	0.22	<b>-2%</b>
	120	240	<b>3053</b>	13.33	<b>+5%</b>
SPONGENT-256	9520	4	<b>1895</b>	0.17	<b>-3%</b>
	140	272	<b>3425</b>	11.43	<b>+4%</b>

# SPONGENT: The Area, NXP 90 nm

Hash Function	Cycles	Datapath [bit]	Area [GE]	Throughput [kb/s]	Area compared to UMC
SPONGENT-88	900	4	<b>532</b>	0.81	<b>-28%</b>
	45	88	<b>892</b>	17.78	<b>-21%</b>
SPONGENT-128	2380	4	<b>748</b>	0.34	<b>-29%</b>
	70	136	<b>1301</b>	11.43	<b>-23%</b>
SPONGENT-160	3960	4	<b>935</b>	0.40	<b>-30%</b>
	90	176	<b>1752</b>	17.78	<b>-20%</b>
SPONGENT-224	7200	4	<b>1206</b>	0.22	<b>-30%</b>
	120	240	<b>2334</b>	13.33	<b>-20%</b>
SPONGENT-256	9520	4	<b>1364</b>	0.17	<b>-30%</b>
	140	272	<b>2612</b>	11.43	<b>-20%</b>



# SPONGENT: Some Security Properties

**Theorem 1.** *Any 5-round differential characteristic of the underlying permutation in SPONGENT- $\{88, 128, 160, 224, 256\}$  has a minimum of 10 active S-boxes.*

# of rounds	SPONGENT-88		SPONGENT-128		SPONGENT-160		SPONGENT-224		SPONGENT-256	
	ASN	Prob	ASN	Prob	ASN	Prob	ASN	Prob	ASN	Prob
5	10	$2^{-21}$	10	$2^{-22}$	10	$2^{-21}$	10	$2^{-21}$	10	$2^{-20}$
10	20	$2^{-47}$	29	$2^{-68}$	20	$2^{-50}$	20	$2^{-43}$	—	—
15	30	$2^{-74}$	-	-	30	$2^{-79}$	30	$2^{-66}$	—	—

Moreover, to thwart linear multidimensional/saturation attacks, the new 4x4-bit S-box is chosen such that the linear hull contains at most one trail with a single active S-box in every round



# Conclusions

- Sponge construction with reduced security claims so far most suitable for low-area hashing
  - almost no overhead in the footprint with respect to the absolute minimum
- The underlying permutation is based on the PRESENT design ideas
  - similar differential and better linear properties
- Multidimensional linear properties (discovered for round-reduced PRESENT) eliminated by the new S-box

# Outlooks I

- SPONGENT is a hermetic design (the permutation is not allowed to have structural weaknesses) → higher numbers of rounds needed
- But this does not seem to be necessary for the concrete (second) preimage and collision security
- **Open problem:** In sponge-type constructions, argue provable security in a more relaxed model (maybe no more indifferenciability) and go for simpler permutations to increase performance in designs

# Outlooks II

- In some applications, only full collision security is necessary

- In permutation-based sponge, one can put

$$\text{Rate} = \frac{1}{2} \text{Capacity}$$

to sacrifice preimage security for speed

→ **Faster SPONGENT variants:** Stay tuned!

- For applications with full 2<sup>nd</sup> preimage security requirements, the state has to be almost doubled

→ **SPONGENT variants with full 2<sup>nd</sup> preimage security:**  
Coming soon!